

1. HANDLING OBELIX

I work with the free multiplicative group G on countably many generators, but an additive (abelian) free group also works. The functional equation for Obelix is

$$f(f^2(h)f(h)^{-1}) = hf(h)^{-1}.$$

Think of this as saying that if $(a, b), (b, c) \in f$, then we must have $(cb^{-1}, ab^{-1}) \in f$.

Define \mathcal{E} as the collection of sets $E \subseteq G^2$ satisfying the following properties.

- (1) E is finite.
- (2) E is an injective function.
- (3) $(1, 1) \in E$.
- (4) If $(a, b), (b, c) \in E$, then $(cb^{-1}, ab^{-1}) \in E$.
- (5) If $(a, b), (c, d) \in E$ and $ab^{-1} = cd^{-1}$, then $a = c$.
- (6) If $(a, b) \in E$ and $b \notin \text{dom}(E)$, then $ab^{-1} \notin \text{dom}(E) \cup \text{im}(E)$.

Lemma 1.1. *For any $E \in \mathcal{E}$ and any $a \in G$, there is an extension $E \subseteq E' \in \mathcal{E}$ where the functional equation holds for a .*

Proof. **Case 1:** Assume $(a, b) \in E$ for some $b \in G$.

If $b \in \text{dom}(E)$, then by condition (4) we are already done. So reduce to the case when $b \notin \text{dom}(E)$. In particular, by (2) and (3) we know $a, b \neq 1$, and also by (6) we know that $ab^{-1} \notin \text{dom}(E) \cup \text{im}(E)$ (and is not 1).

Take c to be a generator of G not appearing in the reduced form of any entry in E , and set $E' := E \cup \{(b, c), (cb^{-1}, ab^{-1})\}$. Conditions (1), (3), and (5) are immediate. Condition (2) is also clear, where injectivity needs $ab^{-1} \notin \text{im}(E)$. Condition (4) is also easy to check, using the fact that E is injective, $cb^{-1} \neq c$, and $ab^{-1} \notin \text{dom}(E)$.

Finally, for condition (6), a finite check works. The main case is checking that $ca^{-1} \notin \text{dom}(E') \cup \text{im}(E')$. This is clear since $a \neq 1$ and $a \neq b$ (by condition (5) for E). One should also note that by condition (5) (and the fact that E is a function), there is no pair $(x, y) \in E$ with $(x, y) \neq (a, b)$ and $xy^{-1} = ab^{-1}$, so we don't ruin condition (6) for pairs already in E .

Case 2: Assume $a \notin \text{dom}(E)$. If $(x, a) \in E$ for some (unique by (2)) $x \in G$, then applying Case 1 to x , we reduce to the case when $a \in \text{dom}(E)$.

Thus, we may consider the case when $a \notin \text{dom}(E) \cup \text{im}(E)$. If $(x, y) \in E$ with $xy^{-1} = a$, then by applying Case 1 to x , we get $a \in \text{im}(E)$, and reduce to a previously considered case. So, we may assume there is no such pair (x, y) . Fixing b to be a generator of G not appearing in the reduced forms for the entries in E , nor in a , then after passing to $E \cup \{(a, b)\} \in \mathcal{E}$ we again reduce to Case 1. \square

The functional equation for Asterix is $f(f(f(h)h^{-1})h) = h$. Taking the seed

$$\{(1, 1), (x_1, x_2), (x_2x_1^{-1}, x_3), (x_3x_1, x_4)\} \in \mathcal{E}$$

works to contradict this equation.