1. Imports and Opened Namespaces

```
import Lean
import Batteries.Tactic.Exact
import Qq
```

open Lean Elab Command Term Meta Qq

- Imports: The code imports necessary modules:
 - Lean: The core Lean language functionalities.
 - Batteries.Tactic.Exact: Likely a custom or additional module providing enhanced tactics (assuming it's available in the environment).
 - Qq: A macro and metaprogramming utility for working with quoted expressions in Lean.
- **Namespaces**: Opens several Lean namespaces to make functions and types available without qualifying them:
 - Lean, Elab, Command, Term, Meta, Qq.

2. Theorems for Commutativity of Equality and Inequality

```
theorem Eq.comm'.{u} {\alpha : Sort u} {a b : \alpha} : (a = b) = (b = a) := propext Eq.comm
theorem Ne.comm'.{u} {\alpha : Sort u} {a b : \alpha} : (a \neq b) = (b \neq a) := propext ne_comm
```

- **Purpose**: These theorems express the commutativity of equality (=) and inequality (≠) in terms of propositional equality (=) rather than logical equivalence (↔).
- Explanation:
 - Eq. comm states that $a = b \leftrightarrow b = a$.
 - propext converts a logical equivalence into a propositional equality.
 - Thus, Eq. comm' states (a = b) = (b = a).

3. Definition of Lean.MVarId.congrWith

```
partial def Lean.MVarId.congrWith (mvarId_ : MVarId) (eqs : List Expr)
: MetaM Unit := ...
```

- **Purpose**: Attempts to solve a metavariable (an unsolved goal in the proof) using congruence rules and provided equations, mirroring how Vampire uses congruence closure in its proofs.
- Explanation:
 - Parameters:
 - mvarId_: The metavariable (goal) to solve.
 - eqs: A list of equations that can be used in the congruence process.
 - Process:
 - Tries to solve the goal using refl (reflexivity).
 - If that fails, it tries to directly assign the goal to one of the provided equations (eqs).
 - If still unsuccessful, it breaks down the goal using congruence (congrCore) and recursively tries to solve the subgoals.
 - If necessary, it attempts to swap the sides of equalities or inequalities using Eq.comm' or Ne.comm' and tries again.

4. Definition of Lean. Expr. eqswap

def Lean.Expr.eqswap (e : Expr) : Expr := ...

- **Purpose**: Swaps the left-hand side (LHS) and right-hand side (RHS) of an equality or inequality expression.
- Explanation:
 - Checks if e is an equality (a = b) or inequality $(a \neq b)$.
 - If so, constructs a new expression with the LHS and RHS swapped (b = a or b \neq a).
 - Returns the original expression if it's neither.

5. Definition of mkEqNeSymm

```
def mkEqNeSymm (h : Expr) : MetaM Expr := ...
```

- **Purpose**: Generates the symmetric version of an equality or inequality proof.
- Explanation:
 - \circ If h is a proof of a = b, returns a proof of b = a using Eq.symm.
 - If h is a proof of a \neq b, returns a proof of b \neq a using Ne.symm.
 - Throws an error if h is neither.

6. Elaboration of superpose

- **Purpose**: Implements the superposition inference rule, which is a key rule in Vampire for handling equalities.
- Explanation:
 - Process:
 - Elaborates (elabTerm) and infers types for eq1 and eq2.
 - Decomposes their types to instantiate any universal quantifiers.
 - Constructs a fresh metavariable v of type e2 = typ, where e2 is the instantiated RHS of eq2 and typ is the target type.
 - Uses congrWith to try to solve v using eq1 and its symmetric version.
 - Constructs the final expression by applying the proven equality to eq2.
 - Ensures that the resulting expression has the correct type and that all metavariables are resolved (using assumption).

7. Definitions for Subsumption and Resolution

Subsumption Functions:

```
partial def doSubsumptionSingle {a : Type} (t1 : Q(Prop)) (target :
Q(Prop)) ...
partial def doSubsumptionWith {a : Type} (res : Q(Prop)) (t2 :
Q(Prop)) (target : Q(Prop)) ...
```

•

- **Purpose**: These functions attempt to determine if one clause (logical formula) subsumes another, meaning one implies the other.
- Explanation:
 - doSubsumptionSingle: Checks if a single term t1 subsumes target, considering commutativity for equalities and inequalities.
 - doSubsumptionWith: Handles the case when t2 is a disjunction (logical OR), recursively checking subsumption for each disjunct.

Resolution Functions:

```
partial def doResolutionSingle {a : Type} (t1 : Q(Prop)) (t2 :
Q(Prop)) ...
partial def doResolution (t1 : Q(Prop)) (t2 : Q(Prop)) ...
```

•

• **Purpose**: Implement the resolution inference rule, which is fundamental in automated theorem proving for propositional and first-order logic.

- Explanation:
 - doResolutionSingle: Attempts to resolve a single term t1 with t2, handling negations and potential commutativity.
 - doResolution: Recursively applies doResolutionSingle to disjunctive terms (logical ORs) in t1.

8. Elaboration of resolve

```
elab "resolve " eq1:term:arg eq2:term:arg : term <= typ => do ...
```

- **Purpose**: Provides a tactic to perform resolution between two clauses, similar to how Vampire applies resolution in its proofs.
- Explanation:
 - Process:
 - Elaborates eq1 and eq2, infers their types, and instantiates any variables.
 - Attempts to resolve eq1 with eq2 to produce a term of type typ.
 - If the initial attempt fails, swaps eq1 and eq2 and tries again.
 - Uses the doResolution function to handle the resolution logic.
 - Applies the resulting function to the appropriate proof (eq1 or eq2).

9. Elaboration of subsumption

elab "subsumption " eq1:term:arg eq2:term:arg : tactic => do ...

- **Purpose**: Implements a tactic to prove that one clause subsumes another, effectively showing that one implies the other.
- Explanation:
 - Process:
 - Retrieves the current goal (goal) and its type (goalType).
 - Attempts to match eq1 or eq2 to the goal, potentially swapping sides of equalities or inequalities.
 - Tries to resolve any remaining metavariables using assumption.
 - Assigns the proof to the goal if successful.

10. Macro Definition for mod_symm

```
macro "mod_symm " ar:term:arg : term => `(by first | exact $ar | exact
Eq.symm $ar | exact Ne.symm $ar)"
```

- **Purpose**: Provides a convenient way to attempt to apply a proof ar or its symmetric version to the current goal.
- Explanation:
 - The macro tries exact \$ar, exact Eq.symm \$ar, and exact Ne.symm \$ar in order.
 - This is useful when the proof might require swapping sides of an equality or inequality.

11. Converting Vampire Proofs into Lean Proofs

• Overall Strategy:

- 1. The code defines functions and tactics that correspond to the inference rules used by Vampire.
- 2. By providing equivalents of Vampire's superposition, resolution, and subsumption in Lean, we can replicate the steps of a Vampire proof within Lean's proof environment.

• Step-by-Step Conversion:

- 1. **Identify Proof Steps**: Break down the Vampire proof into individual inference steps (e.g., applications of superposition, resolution, subsumption).
- 2. **Map to Lean Functions**: For each inference step, use the corresponding Lean function or tactic:
 - Use superpose for superposition steps.
 - Use resolve for resolution steps.
 - Use subsumption for subsumption steps.
- 3. **Elaborate Terms**: Use elabTerm to elaborate terms and instantiate variables, aligning the Vampire proof terms with Lean's representation.
- 4. **Handle Equations and Inequalities**: Use Eq.comm', Ne.comm', and functions like mkEqNeSymm to manage commutativity and symmetry as needed.
- 5. **Solve Metavariables**: Use congrWith and other helper functions to resolve metavariables by matching them with known equations or applying congruence.
- 6. **Construct Proofs**: Assemble the individual proof steps into a complete proof in Lean, ensuring that all types align and that the final proof matches the goal.
- 7. **Verify and Optimize**: Use Lean's type checking and metavariable resolution to verify the correctness of each step and optimize the proof.

Example Workflow

Suppose Vampire has proven a theorem using the following steps:

- 1. Superposition between Equations: From a = b and b = c, derive a = c.
- 2. **Resolution with Negated Clause**: From $\neg a = c$ and a = c, derive a contradiction.

3. **Subsumption**: Show that a clause subsumes another.

Conversion into Lean Proof:

```
Step 1: Use superpose to derive a = c:
lean
Code kopieren
have h1 : a = b := ...
have h2 : b = c := ...
have h3 : a = c := superpose h1 h2
```

Step 2: Use resolve to derive a contradiction:

lean Code kopieren have h4 : $\neg a = c := ...$ have h5 : False := resolve h4 h3

Step 3: Use subsumption to show clause subsumption: lean Code kopieren subsumption h_some_clause h_another_clause