### Towards a formal proof of the Freyd-Mitchell embedding theorem

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### Theorem (Freyd-Mitchell 1964)

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- "Diagram chases over modules are valid in all abelian categories"
- Can even define morphisms this way!

# Is this useful? (1)

#### Peter Johnstone (1977):

Incidentally, the Freyd-Mitchell embedding theorem is frequently regarded as a culmination rather than a starting-point; this is because of what seems to me a misinterpretation (or at least an inversion) of its true significance.

[...] I believe its true import is "If you want to prove something about categories of modules, you might as well work in a general abelian category"—for the embedding theorem ensures that your result will be true in this generality, and by forgetting the explicit structure of module categories you will be forced to concentrate on the essential aspects of the problem.

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Formalization experience:

It might not be strictly necessary, but it's certainly useful!

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- Pseudoelements (Mac Lane 1971, Borceux 1994, mathlib: 2020)
  - Associate to an object X a set of *pseudoelements*, which are equivalence classes of morphisms → X.
  - ▶  $X \rightarrow Y$  induces PseudoElem $(X) \rightarrow$  PseudoElem(Y)
  - ▶  $X \to Y$  epi iff PseudoElem $(X) \to$  PseudoElem(Y) surjective, etc.

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```
mono_of_zero_of_map_zero _ fun c hc =>
  have : h c = 0 :=
    suffices \delta (h c) = 0 from zero_of_map_zero _ (pseudo_injective_of_mono _) _ this
    calc
      \delta (h c) = h' (y c) := by rw [\leftarrow Pseudoelement.comp_apply, \leftarrow comm<sub>3</sub>, Pseudoelement.comp_apply]
      _ = h' 0 := by rw [hc]
      _ = 0 := apply_zero _
  Exists.elim ((pseudo_exact_of_exact hgh).2 _ this) fun b hb =>
    have : q'(\beta b) = 0 :=
      calc
        g' (\beta b) = y (g b) := by rw [- Pseudoelement.comp_apply, comm<sub>2</sub>, Pseudoelement.comp_apply]
        = y c := by rw [hb]
        = 0 := hc
    Exists.elim ((pseudo_exact_of_exact hf'g').2 _ this) fun a' ha' =>
      Exists.elim (pseudo_surjective_of_epi \alpha a') fun a ha =>
        have : f a = b :=
           suffices \beta (f a) = \beta b from pseudo injective of mono this
           calc
```

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Problem: the embedding theorem is a lot of work to prove! Pseudoelements (Mac Lane 1971, Borceux 1994, mathlib: 2020) Refinements (Bergman 1974, mathlib: Riou 2023) apply mono\_of\_cancel\_zero intro A  $f_2$  h<sub>1</sub> have  $h_2$  :  $f_2 \ge R_1$ .map' 2 3 = 0 := by rw [ $\leftarrow$  cancel\_mono (app'  $\phi$  3 \_), assoc, NatTrans.naturality, reassoc\_of% h<sub>1</sub>, zero\_comp, zero\_comp] obtain  $\langle A_1, \pi_1, \dots, f_1, hf_1 \rangle$  := (hR<sub>1</sub>'.exact 0).exact\_up\_to\_refinements f<sub>2</sub> h<sub>2</sub> dsimp at hf<sub>1</sub> have  $h_3$  : (f<sub>1</sub> > app'  $\phi$  1) > R<sub>2</sub>.map' 1 2 = 0 := by rw [assoc,  $\leftarrow$  NatTrans.naturality,  $\leftarrow$  reassoc\_of% hf<sub>1</sub>, h<sub>1</sub>, comp\_zero] obtain  $\langle A_2, \pi_2, \dots, g_0, hg_0 \rangle$  := (hR<sub>2</sub>.exact 0).exact\_up\_to\_refinements \_ h<sub>3</sub> obtain  $\langle A_3, \pi_3, \dots, f_0, hf_0 \rangle$  := surjective\_up\_to\_refinements\_of\_epi (app'  $\varphi 0$  \_)  $g_0$ have  $h_4$  :  $f_0 \gg R_1$ .map'  $0 = \pi_3 \gg \pi_2 \gg f_1 := by$ rw [ $\leftarrow$  cancel\_mono (app'  $\phi$  1 \_), assoc, assoc, assoc, NatTrans.naturality,  $\leftarrow$  reassoc\_of% hf<sub>0</sub>, hg<sub>0</sub>]

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$$egin{aligned} Y(P) &: \mathcal{C} o \mathsf{Mod}_R \ X &\mapsto \mathsf{Hom}(P,X) \end{aligned}$$

This is

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This is

- always left exact,
- right exact iff P is projective,
- ▶ faithful iff *P* is a generator,
- $\blacktriangleright$  full as long as all of the above hold and  ${\cal C}$  is cocomplete.

Not every small abelian category is cocomplete and has a projective generator!

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- $\blacktriangleright$  Idea: embed  ${\cal C}$  into a category that has these properties.
- ▶ Freyd uses  $Y : C \to Lex(C^{op}, Ab)$ .
- ▶ Kashiwara-Schapira (and we) use  $Y : C \to Pro(C) \cong Ind(C^{op})^{op} \cong Lex(C^{op}, Ab)$ .
- ▶ Pro(C) works because it is a co-Grothendieck category.

Rough list of things to do to prove the theorem:

- Define Grothendieck categories (easy enough)
- Show that Grothendieck categories are complete (SAFT, done)
- Show that Grothendieck categories have an injective cogenerator (hard)
- Define Pro(C) and prove basic properties (hard, very long, working on it)
- Show that Pro(C) is a co-Grothendieck category (hopefully doable)
- ▶ Prove various properties of the two embeddings  $C \to Pro(C) \to Mod_R$  (easy but somewhat long, mostly done)

Other proofs are possible.

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- We work on this in our free time on weekends
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- How to combat "quadratic bitrot"?
- Solution: use a proof that develops a lot of general theory and PR everything to mathlib immediately
- 138 PRs merged so far (97 to mathlib3, 61 to mathlib4)
- Amount of unmerged code: < 300 lines</p>
- Additional benefits:
  - Do things right from the start, no accumulating piles of hacks (but cannot take any shortcuts)
  - Partial progress is useful to other people

# Size issues (1)

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#### Theorem

Let C be a small abelian category. Then there is a (not necessarily commutative) ring R and a full, faithful and exact functor  $C \to Mod_R$ .

Every category is small if you look at it from far enough away...

### Theorem (Yoneda 1954)

Let C be a category, let  $F : C \to Set$  be a functor and let X be an object of C. Then there is a bijection between the natural transformations  $Hom(X, -) \to F$  and FX.



Advantage when defining data is much less clear.